

LECTURE 10 HOUSEHOLDER REFLECTORS

QR

Previously : GS amounts to the successive application of upper triangular matrices R_i :

$$A \underbrace{R_1 R_2 R_3 \dots R_n}_{R^{-1} \text{ (upper triangular)}} = Q$$

An alternative QR : Householder. We create Q_i 's so that

$$\underbrace{Q_n Q_{n-1} \dots Q_2 Q_1}_{Q^*} A = R, \text{ since}$$

Q_i 's are unitary $\Rightarrow Q$ is unitary.

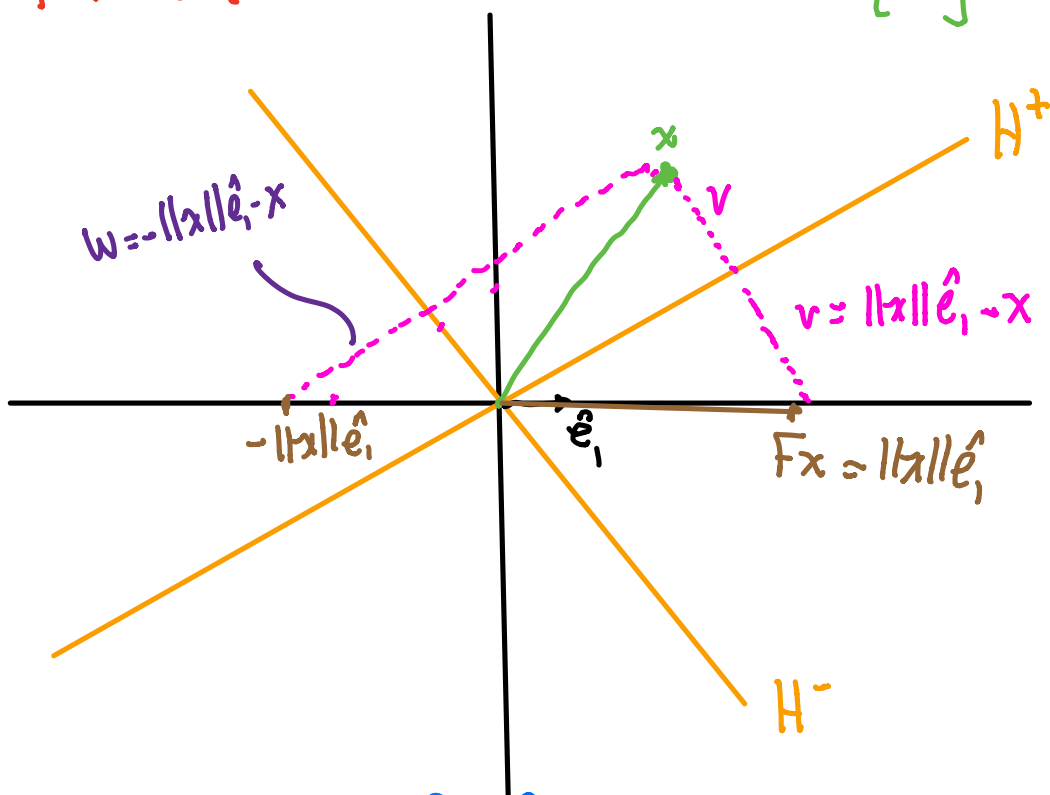
$$\begin{array}{c} \left[\begin{array}{c|c|c} | & | & | \\ \hline & A & \\ \hline | & | & | \end{array} \right] \xrightarrow{Q_1} \underbrace{\left[\begin{array}{c|c|c} x & | & | \\ \hline 0 & | & | \\ \hline \vdots & | & | \end{array} \right]}_{Q_1 A} \xrightarrow{Q_2} \underbrace{\left[\begin{array}{c|c|c} x & x & | \\ \hline 0 & x & | \\ \hline | & 0 & | \end{array} \right]}_{Q_2 Q_1 A} \xrightarrow{Q_3} \underbrace{\left[\begin{array}{c|c|c} x & x & x \\ \hline & x & x \\ \hline & 0 & x \\ & & 0 \end{array} \right]}_{Q_3 Q_2 Q_1 A} = R \end{array}$$

Q_k introduces zeros below the k^{th} row in the k^{th} column w/o destroying the zeros from before

Each

$$Q_k = \begin{bmatrix} \overset{\leftarrow k-1}{I} & \overset{\uparrow k-1}{0} \\ 0 & \underset{\leftarrow m-k+1}{F} \end{bmatrix} \overset{\uparrow m-k+1}{} \quad \overset{e_1}{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \end{bmatrix}$$

F is the Householder Reflector



$$x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \end{bmatrix} \xrightarrow{F} \begin{bmatrix} \|x\| \\ 0 \\ \vdots \end{bmatrix} = \|x\| \hat{e}_1$$

To find the reflector about $H^\perp$:

$$Py = y - \frac{vv^*}{v^*v} y = y - \left(\frac{v^*y}{v^*v} \right) v$$

for some $y \in \mathbb{C}^m$

is the projection of some vector y onto $H^\perp$.

To get to the $\hat{e}_1$ axis:

$$Fy = \left(I - \frac{2vv^*}{v^*v} \right) y = y - 2v \left(\frac{v^*y}{v^*v} \right)$$

$$\therefore \text{Matrix } F = I - \frac{2vv^*}{v^*v}$$

Rank: the projector P (rank $n-1$) and the reflector F (full rank & unitary) only differ by the presence of the 2 !

The Better of 2 Reflectors:

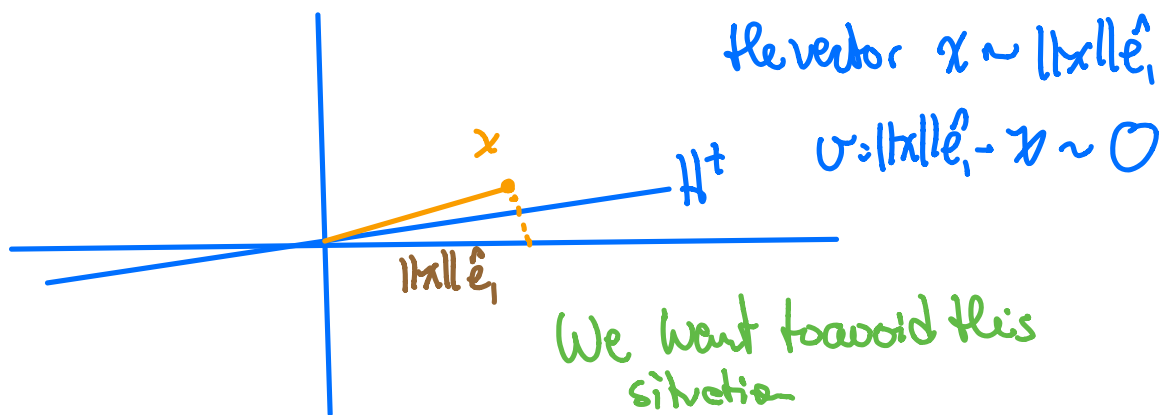
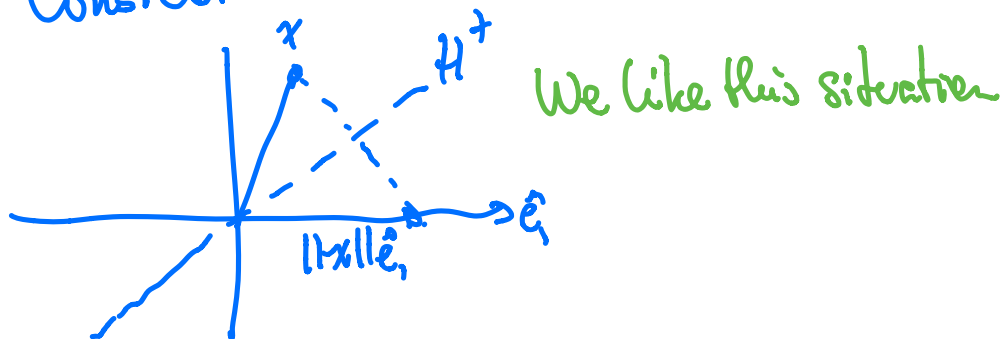
The vector $x \xrightarrow{F} z$ $\|x\| = \|z\|$, $z \in \mathbb{C}$ w/ $|z|=1$

$\therefore$ a circle of possible reflections

If $z \in \mathbb{R}$ then $z = \pm 1$ lead to H^+ & H^- cases,
either is OK, but choice is numerically driven.

We want to avoid subtracting nearby quantities.

Consider



$\therefore$ for this second case why not reflect
about H^- ? (see figure above)

Choose $v = -\text{sgn}(x_1)||x||\hat{e}_1 - x$
(if $x_1 = 0 \Rightarrow \text{sgn} = 0$)

then $\|r\|$ is never smaller than $\|x\|$ //

See Algorithm 10.1

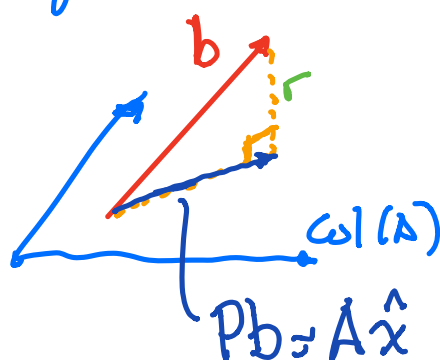
Householder QR

Cost $\sim O(mn^2) \sim O(n^3)$ if square.

but more stable numerically. //

Lecture 11 LSQ (least squares)

Solving $Ax=b$



$$r = Ax - b$$

if $r=0 \Rightarrow Ax=b$

$$Ax = Pb, \quad b = Pb + r$$

Choose r to be "shortest"

$$\|r\|^2 = \|b - Pb\|^2 = \|(I-P)b\|^2$$

the minimizer will be $\hat{x}$. $r \perp \text{col}(A)$

The orthogonality condition

$$A^* r = 0$$

$$\text{or } A^*(b - A\hat{x}) = 0$$

$$\text{or } A^* A \hat{x} = A^* b \quad \text{The Normal Equations}$$

$$\text{solving for } \hat{x} = (A^* A)^{-1} A^* b$$

$$\text{which satisfies } A\hat{x} = Pb = A(A^* A)^{-1} A^* b$$

$\underbrace{\hspace{10em}}_P$

//

$$\text{ex) } Ax = b \quad A = \begin{bmatrix} 1 & 2 \\ 1 & 3 \\ 0 & 0 \end{bmatrix} \quad b = \begin{bmatrix} 4 \\ 5 \\ 6 \end{bmatrix}$$

" "
 $a_1 \quad a_2$

so clear that $b \neq x_1 a_1 + x_2 a_2$

but $Pb = \hat{x}_1 a_1 + \hat{x}_2 a_2$ can be found.

$$\text{so let } r = b - Pb = (I - P)b$$

Find $\hat{x}$ s.t. $A^T r = 0$

$A^T A \hat{x} = A^T b$ the Normal Eqs

$$A^T A = \begin{bmatrix} 1 & 1 & 0 \\ 2 & 3 & 0 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 1 & 3 \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} 2 & 5 \\ 5 & 13 \end{bmatrix}$$

$$\hat{x} = (A^T A)^{-1} A^T b = \begin{bmatrix} 13 & -5 \\ -5 & 2 \end{bmatrix} \begin{bmatrix} 1 & 1 & 0 \\ 2 & 3 & 0 \end{bmatrix} \begin{bmatrix} 4 \\ 5 \\ 6 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

$$Pb = A \hat{x} = \begin{bmatrix} 1 & 2 \\ 1 & 3 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 4 \\ 5 \\ 0 \end{bmatrix}$$

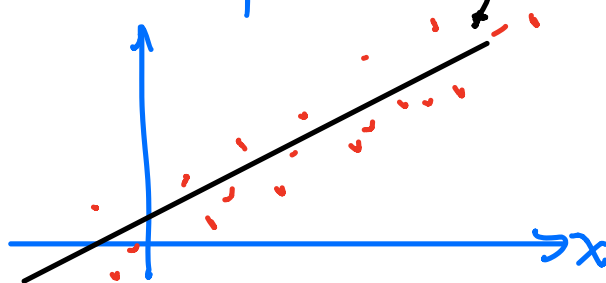
$$b - Pb = r = \begin{bmatrix} 0 \\ 0 \\ 6 \end{bmatrix} = (I - P)b \quad \therefore \|r\|_2 = 6 //$$

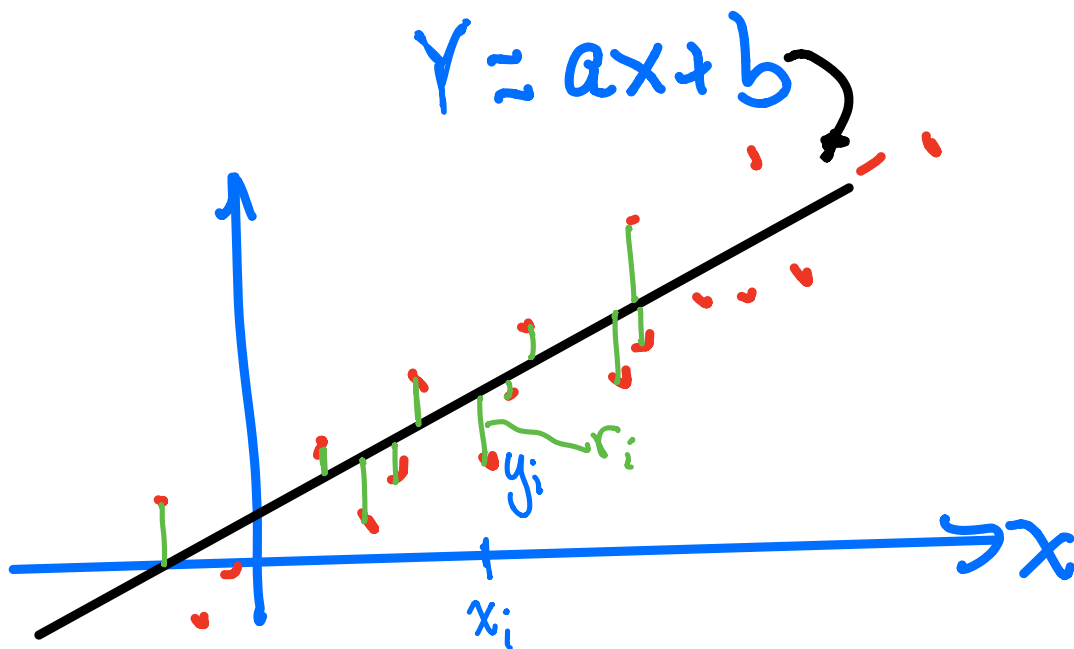
FITTING DATA TO MODELS

Data:

x_1	y_1
x_2	y_2
$\vdots$	$\vdots$
$\vdots$	$\vdots$
x_m	y_m

"Model" look like a straight line
 $y = ax + b$





$$r_i = Y(x_i) - y_i$$

Variance
Energy = $\sum |r_i|^2$

Find a, b s.t. energy is minimized.

Least squares problem:

$$\text{let } A = \begin{bmatrix} 1 & x_1 \\ 1 & x_2 \\ 1 & x_3 \\ \vdots & \vdots \\ 1 & x_m \end{bmatrix} \quad Z = \begin{bmatrix} a \\ b \end{bmatrix} \quad y = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}$$

$$r = \begin{bmatrix} r_1 \\ r_2 \\ \vdots \\ r_m \end{bmatrix}$$

$$AZ = y + r$$

Solve for $\hat{Z}$ s.t. $\|r\|_2^2$ is smallest.

POLYNOMIAL DATA FITTING

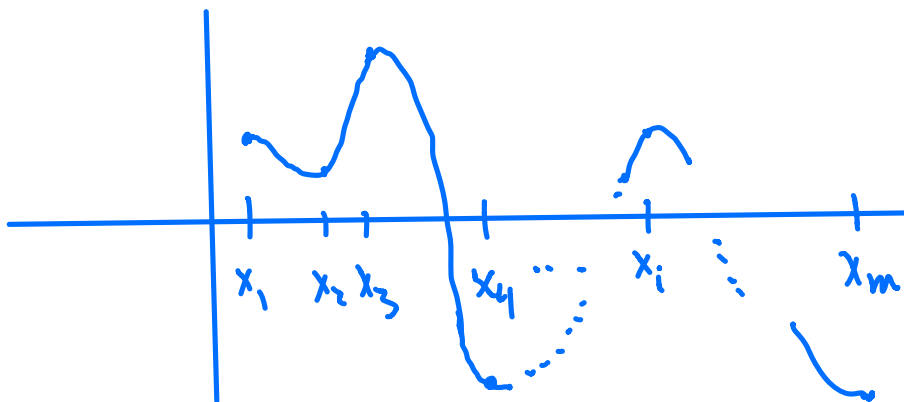
$$\{x_i\}_{i=1}^m, \{y_i\}_{i=1}^m \quad \begin{matrix} x \in \mathbb{C}^m \\ y \in \mathbb{C}^m \end{matrix}$$

$\exists$! polynomial interpolant

$$p(x) = c_0 + c_1x + c_2x^2 + \dots + c_{m-1}x^{m-1}$$

$$\text{s.t. at } x_i \quad p(x_i) = y_i$$

m eqs & m unknowns, the powers of x are linearly independent.



The condition $p(x_i) = y_i \quad i=1, 2, \dots, m$
generates

Remark: in interpolation we insist that $p(x_i) = y_i, i=1, \dots, m$
 $\therefore$ we need m equations for m unknowns.

$$\begin{bmatrix} 1 & x_1 & x_1^2 & \dots & x_1^{n-1} \\ 1 & x_2 & x_2^2 & \dots & x_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_m & x_m^2 & \dots & x_m^{n-1} \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_{n-1} \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}$$

↑
Vandermonde
Matrix

So imagine fitting 10^6 points... using
a polynomial of deg $10^6 - 1$??

... Highly unstable process on a finite
precision machine.

Remark: in approximation we don't insist that $p(x_i) = y_i$
for $i = 1, \dots, m$. Instead we propose an approximating
polynomial that has some optimizing
characteristic. For example, we can insist
that we find an approximating polynomial
that minimizes a residual $p(x_i) - y_i$ //