
Random Self-Similar Trees: Dynamical Pruning, Invariance, and Criticality

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Trees.

$\mathcal{L}_{\text{plane}}$ - space of finite unlabeled rooted reduced binary trees with edge lengths and planar embedding.

The space $\mathcal{L}_{\text{plane}}$ includes the empty tree $\phi = \{\rho\}$ comprised of a root vertex ρ and no edges.

$d(x, y)$: the length of the minimal path within T between x and y .

The length of a tree T is the sum of the lengths of its edges:

$$\text{length}(T) = \sum_{i=1}^{\#T} l_i.$$

The height of a tree T is the maximal distance between the root and a vertex:

$$\text{height}(T) = \max_{1 \leq i \leq \#T} d(v_i, \rho).$$

Partial ordering.

Consider a tree $T \in \mathcal{L}_{\text{plane}}$ and a point $x \in T$. Let $\Delta_{x,T}$ denote all points of T descendant to x , including x . Then $\Delta_{x,T}$ is itself a tree in $\mathcal{L}_{\text{plane}}$ with root at x .

Let (T_1, d) and (T_2, d) be two metric rooted trees, and let ρ_1 denote the root of T_1 . A function

$$f : (T_1, d) \rightarrow (T_2, d)$$

is an **isometry** if

$$\text{Image}[f] \subseteq \Delta_{f(\rho_1), T_2}$$

and $\forall x, y \in T_1$,

$$d(f(x), f(y)) = d(x, y).$$

Partial order: $T_1 \preceq T_2$ if and only if \exists an isometry $f : (T_1, d) \rightarrow (T_2, d)$.

Generalized dynamical pruning.

Consider a **monotone non-decreasing**

$$\varphi : \mathcal{L}_{\text{plane}} \rightarrow \mathbb{R},$$

i.e. $\varphi(T_1) \leq \varphi(T_2)$ whenever $T_1 \preceq T_2$.

Generalized dynamical pruning operator

$$\mathcal{S}_t(\varphi, T) : \mathcal{L}_{\text{plane}} \rightarrow \mathcal{L}_{\text{plane}}$$

induced by φ at any $t \geq 0$:

$$\mathcal{S}_t(\varphi, T) := \rho \cup \left\{ x \in T \setminus \rho : \varphi(\Delta_{x,T}) \geq t \right\}.$$

\mathcal{S}_t cuts all subtrees for which the value of φ is below threshold t . Here,

$$\mathcal{S}_s(T) \preceq \mathcal{S}_t(T)$$

whenever $s \geq t$.

Example: Tree height.

Recall: $\mathcal{S}_t(\varphi, T) := \rho \cup \left\{ x \in T \setminus \rho : \varphi(\Delta_{x,T}) \geq t \right\}$.

Let the function $\varphi(T)$ equal the **height** of T :

$$\varphi(T) = \text{height}(T).$$

Semigroup property: $\mathcal{S}_t \circ \mathcal{S}_s = \mathcal{S}_{t+s}$ for any $t, s \geq 0$.

It coincides with the **tree erasure** introduced by Neveu (1986).

Neveu (1986): established invariance of a critical and sub-critical binary Galton-Watson processes with i.i.d. exponential edge lengths with respect to the tree erasure.

Example: Total tree length.

Recall: $\mathcal{S}_t(\varphi, T) := \rho \cup \left\{ x \in T \setminus \rho : \varphi(\Delta_{x,T}) \geq t \right\}$.

Let the function $\varphi(T)$ equal the **total lengths** of T :

$$\varphi(T) = \text{length}(T).$$

In this case the pruning operator \mathcal{S}_t coincides with the potential dynamics of **Burgers equation**, as shown in

M. Arnold, YK, I. Zaliapin (2017) - arXiv:1707.01984

Example: Horton pruning.

Let

$$\varphi(T) = k(T) - 1,$$

where the **Horton-Strahler order** $k(T)$ is the minimal number of Horton prunings \mathcal{R} (cutting the tree leaves and applying series reduction) necessary to eliminate the tree T .

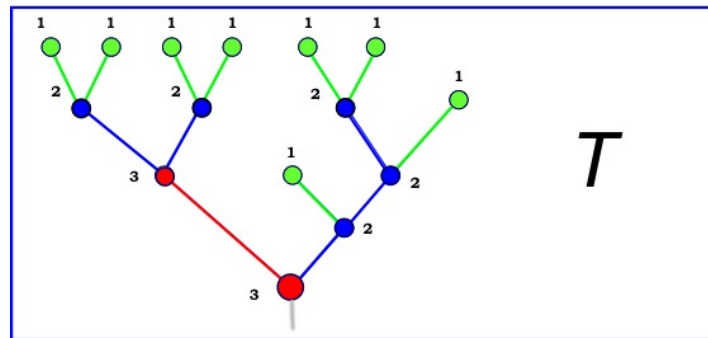
Here,

$$\mathcal{S}_t = \mathcal{R}^{[t]}.$$

The Horton-Strahler order is known as the **register number** as it equals the minimum number of memory registers necessary to evaluate an arithmetic expression described by a tree T .

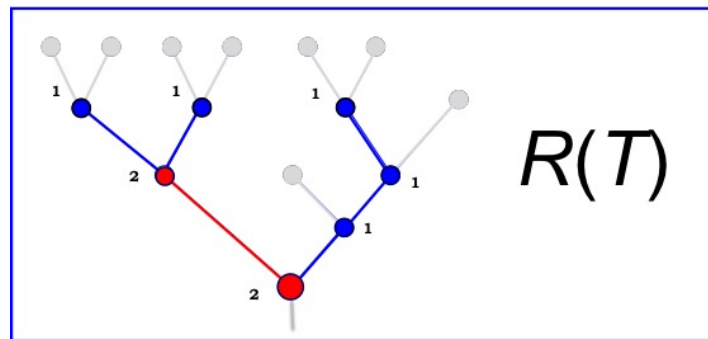
Horton pruning and Horton-Strahler ordering.

- **Pruning** $\mathcal{R}(T)$ of a finite tree T cuts the leaves, followed by *series reduction*.
- A chain of the same order vertices with edges connecting to parent vertices is called **branch**.
- Branches cut at k -th pruning, $\mathcal{R}^{k-1}(T) \setminus \mathcal{R}^k(T)$, have order k , $k \geq 1$.
- N_k denotes the number of **branches** of order k in a finite tree T



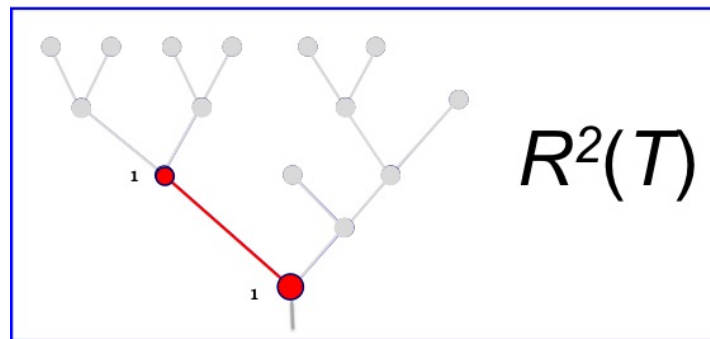
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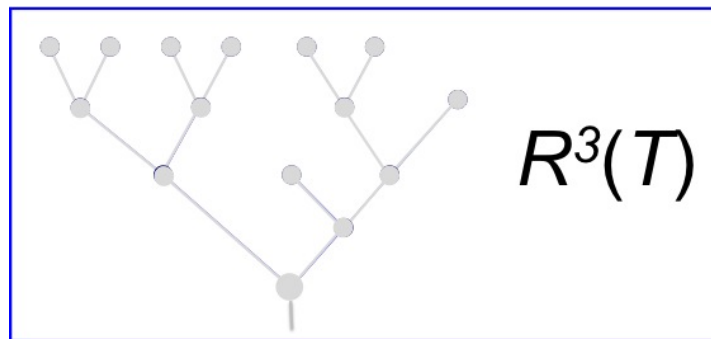
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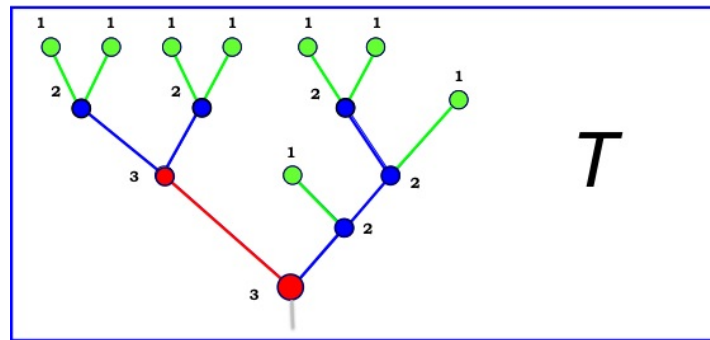
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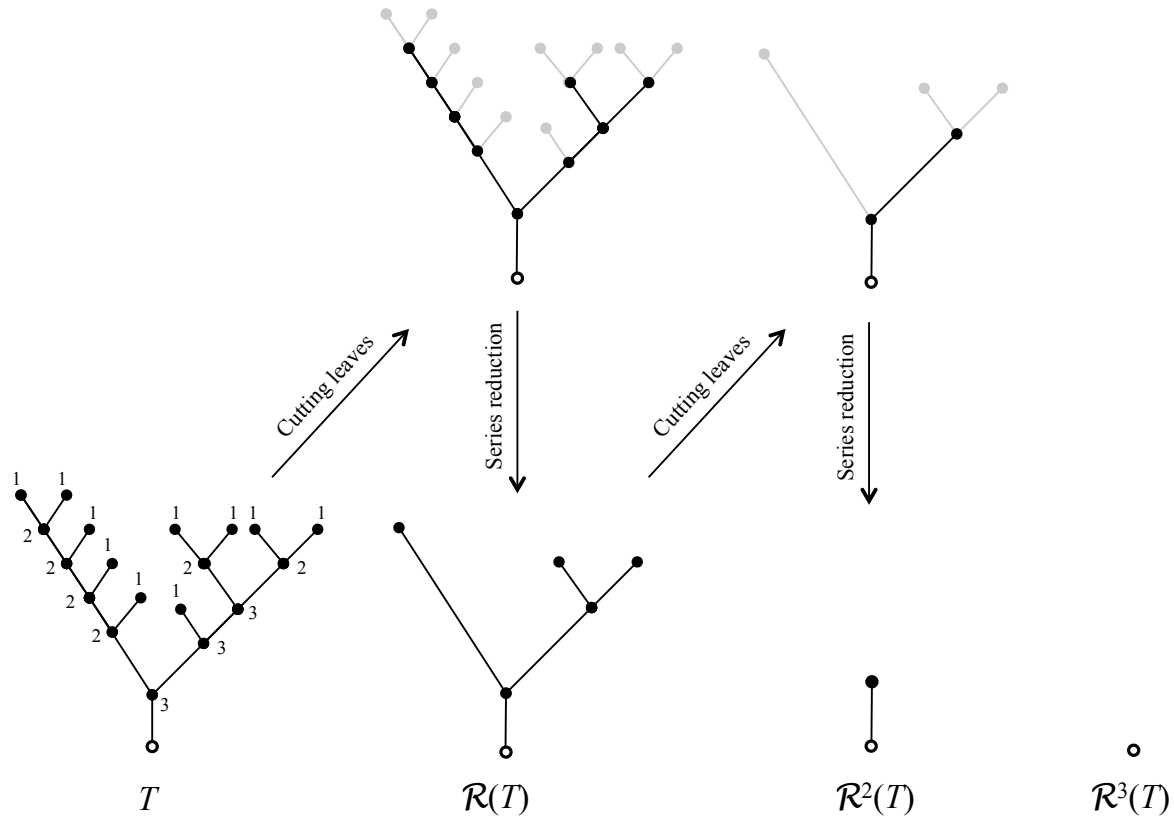


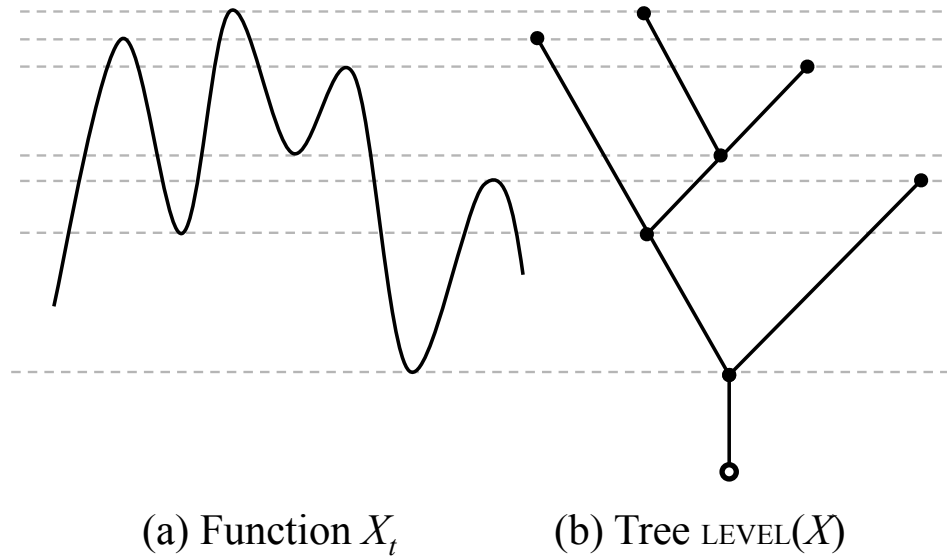
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Pruning of a tree mod series reduction



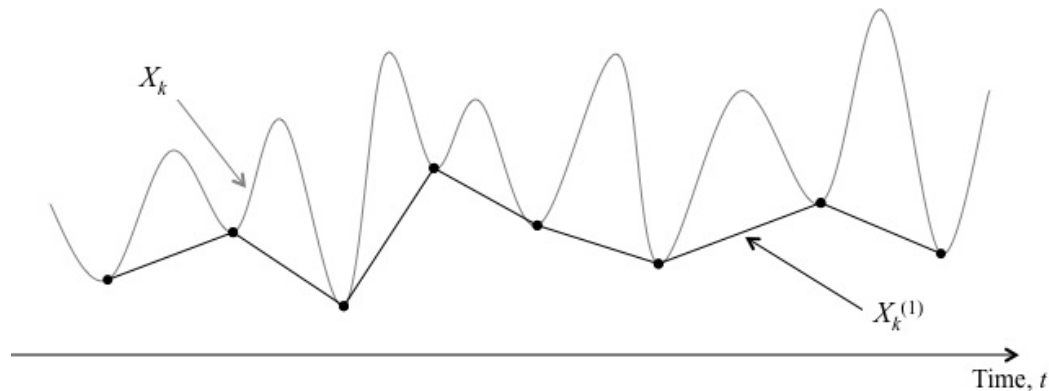
Level-set tree of a function.

Function X_t (panel a) with a finite number of local extrema and its level-set tree $\text{level}(X)$ (panel b).

Horton pruning of time series

The transition from a time series X_k to the time series $X_k^{(1)}$ of its local minima corresponds to [Horton pruning](#) of the level-set tree $\text{level}(X)$.

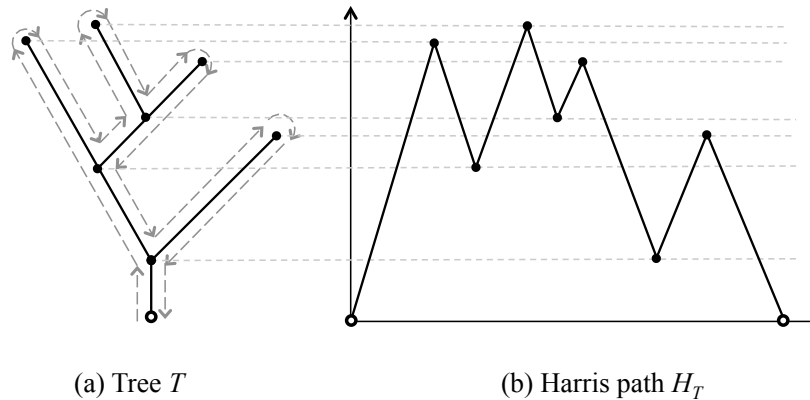
Zaliapin and YK, CSF (2012).



Exponential critical binary Galton-Watson tree

We say that a random tree $T \in \mathcal{L}_{\text{plane}}$ is an **exponential critical binary Galton-Watson tree** with parameter $\lambda > 0$, and write $T \stackrel{d}{=} \text{GW}(\lambda)$, if

- (i) $\text{shape}(T)$ is a critical binary Galton-Watson tree;
- (ii) the orientation for every pair of siblings in T is uniformly random and symmetric;
- (iii) given $\text{shape}(T)$, the edges of T are sampled as independent exponential random variables with parameter λ .

Exponential critical binary Galton-Watson tree

The level set tree $\text{level}(X_t)$ is an exponential critical binary Galton-Watson tree $\text{GW}(\lambda)$ if and only if the rises and falls of X_t , excluding the last fall, are distributed as independent exponential random variables with parameter $\lambda/2$.

J. Neveu and J. Pitman (1989), J. F. Le Gall (1993)

Invariance under pruning

Theorem. [M. Arnold, YK, I. Zaliapin, 2017]

Let $T \stackrel{d}{=} \text{GW}(\lambda)$ be an exponential critical binary Galton-Watson tree with parameter $\lambda > 0$.

Then, for any monotone non-decreasing function $\varphi : \mathcal{L}_{\text{plane}} \rightarrow \mathbb{R}^+$ we have

$$T^t := \{\mathcal{S}_t(\varphi, T) \mid \mathcal{S}_t(\varphi, T) \neq \phi\} \stackrel{d}{=} \text{GW}(\lambda p_t(\lambda, \varphi)),$$

where $p_t(\lambda, \varphi) = \text{P}(\mathcal{S}_t(\varphi, T) \neq \phi)$.

That is, the pruned tree T^t conditioned on surviving is an exponential critical binary Galton-Watson tree with parameter

$$\mathcal{E}_t(\lambda, \varphi) = \lambda p_t(\lambda, \varphi).$$

Invariance under pruning**Theorem.** [M. Arnold, YK, I. Zaliapin, 2017]

(a) If $\varphi(T)$ equals the total length of T ($\varphi = \text{length}(T)$), then

$$\mathcal{E}_t(\lambda, \varphi) = \lambda e^{-\lambda t} [I_0(\lambda t) + I_1(\lambda t)].$$

(b) If $\varphi(T)$ equals the height of T ($\varphi = \text{height}(T)$), then

$$\mathcal{E}_t(\lambda, \varphi) = \frac{2\lambda}{\lambda t + 2}.$$

(c) If $\varphi(T) + 1$ equals the Horton-Strahler order of the tree T , then

$$\mathcal{E}_t(\lambda, \varphi) = \lambda 2^{-\lfloor t \rfloor}.$$

Distributional prune-invariance

Definition. Consider a probability measure μ on $\mathcal{L}_{\text{plane}}$ such that $\mu(\phi) = 0$. Let

$$\nu(T) = \mu \circ \mathcal{S}_t^{-1}(T) = \mu(\mathcal{S}_t^{-1}(T)).$$

Measure μ is called **invariant** with respect to the pruning operator $\mathcal{S}_t(\varphi, T)$ if for any tree $T \in \mathcal{L}_{\text{plane}}$ we have

$$\mu(T) = \nu(T|T \neq \phi).$$

Also need the invariance of the distribution of edge lengths in the pruned tree $T_t := \mathcal{S}_t(\varphi, T)$.

YK and I. Zaliapin (2017) - arXiv:1608.05032

Open question: finding and classifying all the invariant probability measures μ on $\mathcal{L}_{\text{plane}}$.

Other prune-invariances

YK and Zaliapin (Fractals 2016)

- **Coordination** and **mean self-similarity** imply **strong Horton law**.

YK & Zaliapin (AIHP 2017):

- Established the **root-Horton law** for the Kingman's coalescent.
- Showed that the tree for Kingman's coalescent is combinatorially equivalent to the level-set tree of iid time series (the two are **one Horton pruning apart**).

Perspectives.

- Time series: extreme values.
- Generalized notions of self-similarity under Horton pruning: **YK and I. Zaliapin (2017) - arXiv:1608.05032**

A class of multi-type branching processes is considered: **the Hierarchical Branching Processes**

- Generalized dynamical pruning. Burgers equations. **M. Arnold, YK, I. Zaliapin (2017) - arXiv:1707.01984**
- Models of statistical mechanics at criticality.